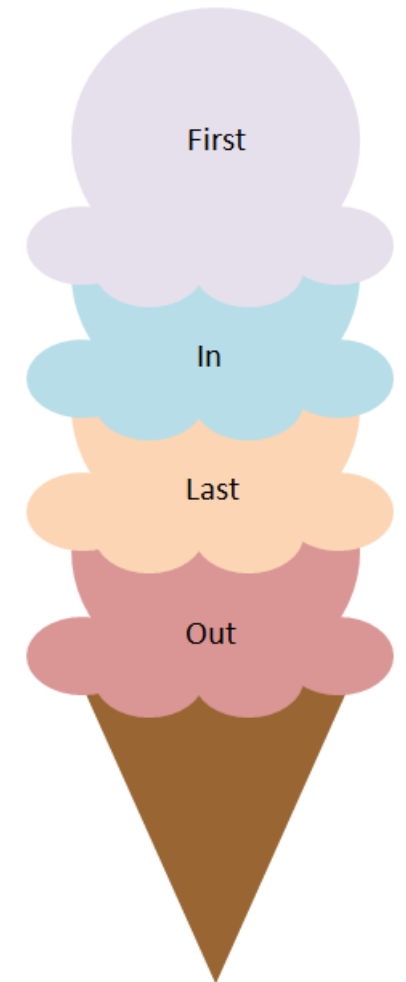


Stack



Stack ADT

- A **stack** is a data structure in which all access is restricted to the most recently inserted element.
 - Stack has only one end.
- Insertions and deletions follow last-in first-out (**LIFO**) scheme (principle).
 - It means the element added last will be removed first.
- **Main operations:**
 - **push(object)**: insert element
 - **object pop()**: remove and returns last element
- **Auxiliary operations:**
 - **object top()**: returns last element without removing it.
 - **integer size()**: returns number of elements stored.
 - **boolean isEmpty()**: returns whether no elements are stored.

Applications of Stacks

➤ Direct

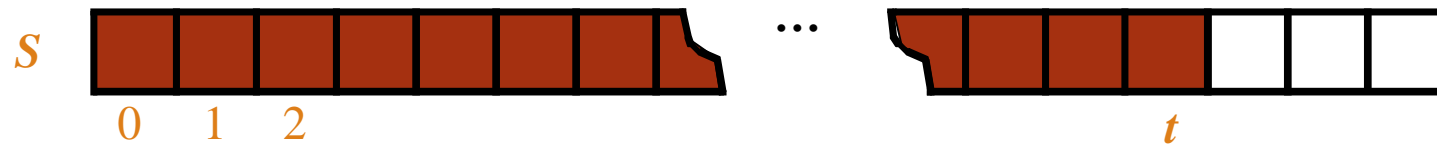
- Page visited history in a web browser.
- Undo sequence in a text editors.
- Chain of method calls in C++ runtime environment.
- Stack is used to evaluate prefix, postfix and infix expressions.
- An expression can be represented in prefix, postfix or infix notation. Stack can be used to convert one form of expression to another.

➤ Indirect

- Auxiliary data structure for algorithms.
- Component of other data structures.

Array-based Stack

- Add elements in an array S of capacity(size) N .
- A variable top keeps track of the index of the top element.
- Size is $top+1$



Push and Pop Algorithms

Algorithm *push(Element)*:

```
if top = N-1 then
    throw “Full Stack Exception”
else
    top ← top + 1
    S[top] ← Element
```

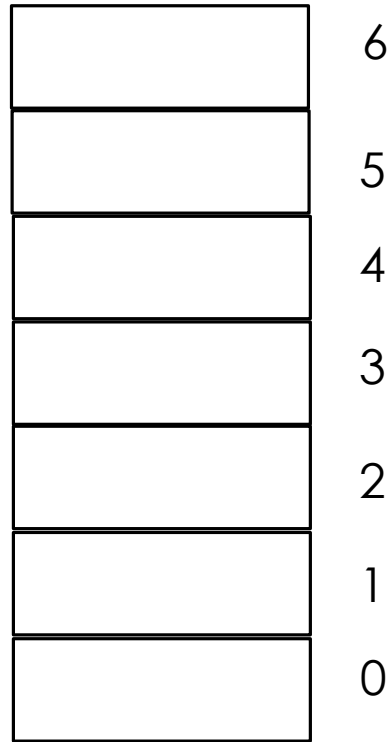
Run time: $O(1)$

Algorithm *pop()*:

```
if isEmpty() then
    throw “Empty Stack Exception”
else
    top ← top - 1
    return S[top + 1]
```

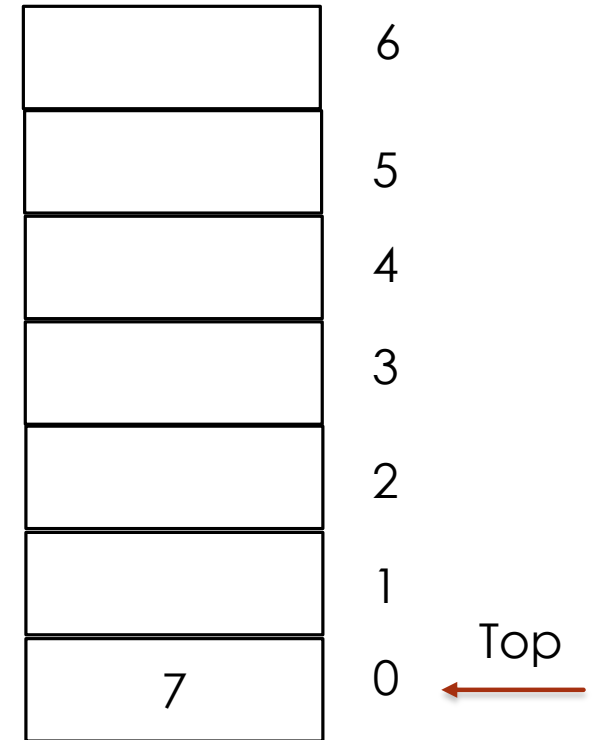
Run Time: $O(1)$

Stack Operations - Example



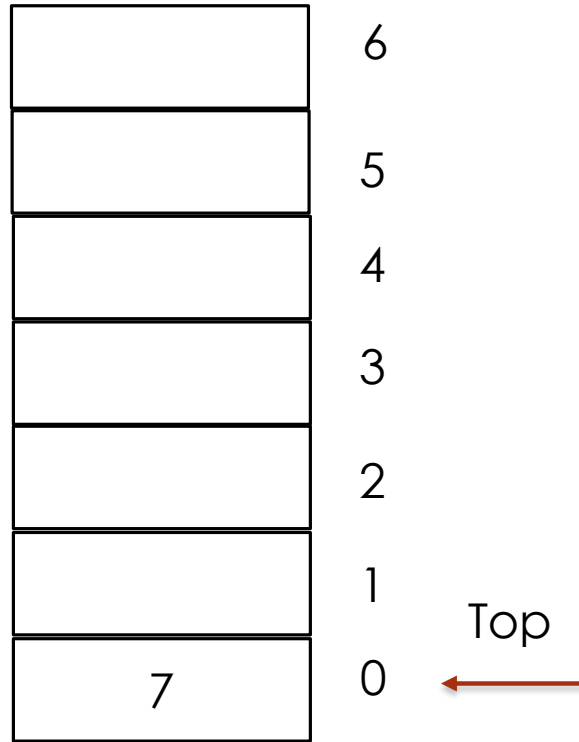
Stack S, N=7, Top=-1

```
Push (7)  
{  
  if top==N-1 Then  
    “Overflow”  
  else  
    Top=Top+1  
    S[top]=7  
}
```

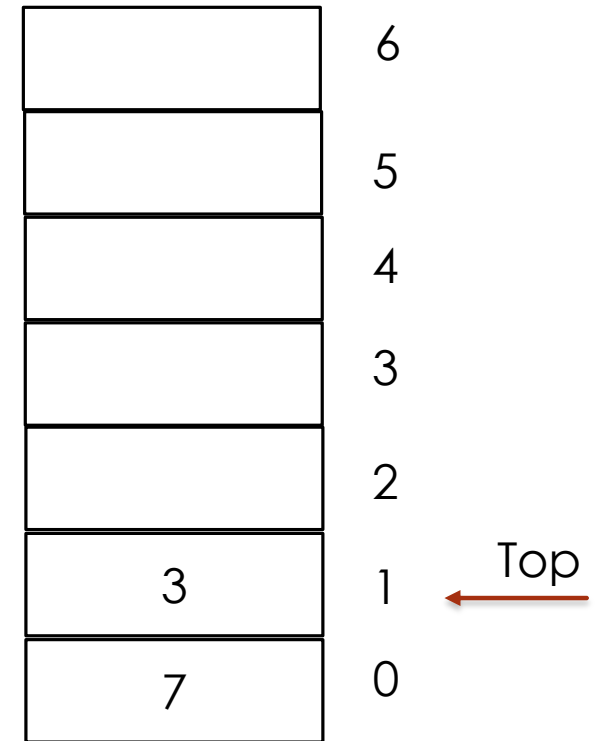


Stack S, N=7

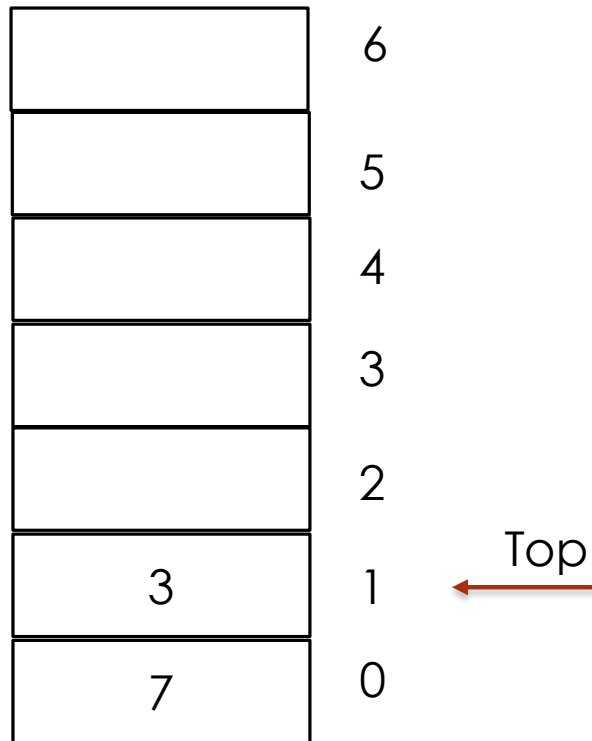
Stack Operations - Example



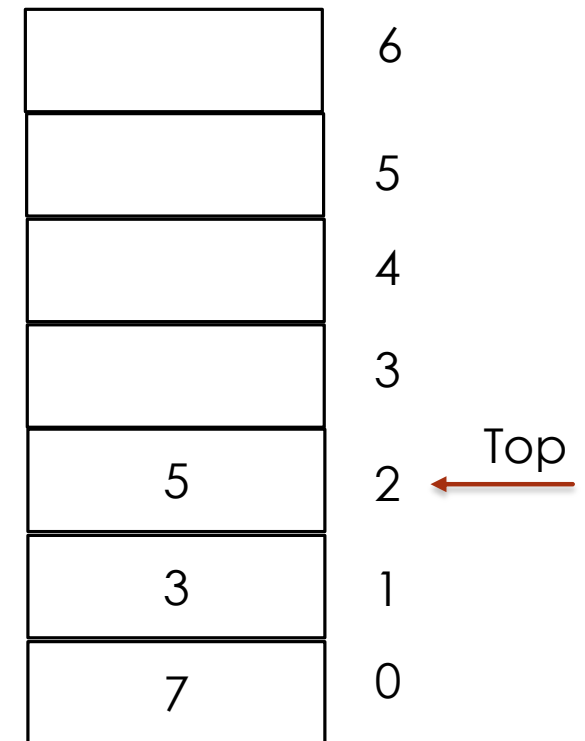
```
Push (3)  
{  
  if top==N-1 Then  
    "Overflow"  
  else  
    Top=Top+1  
    S[top]=3  
}
```



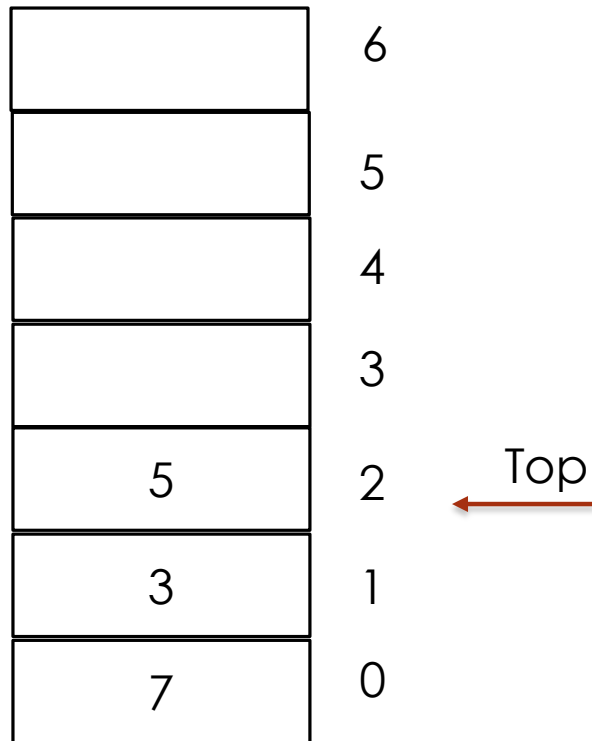
Stack Operations - Example



```
Push (5)  
{  
  if top==N-1 Then  
    “Overflow”  
  else  
    Top=Top+1  
    S[top]=5  
}
```

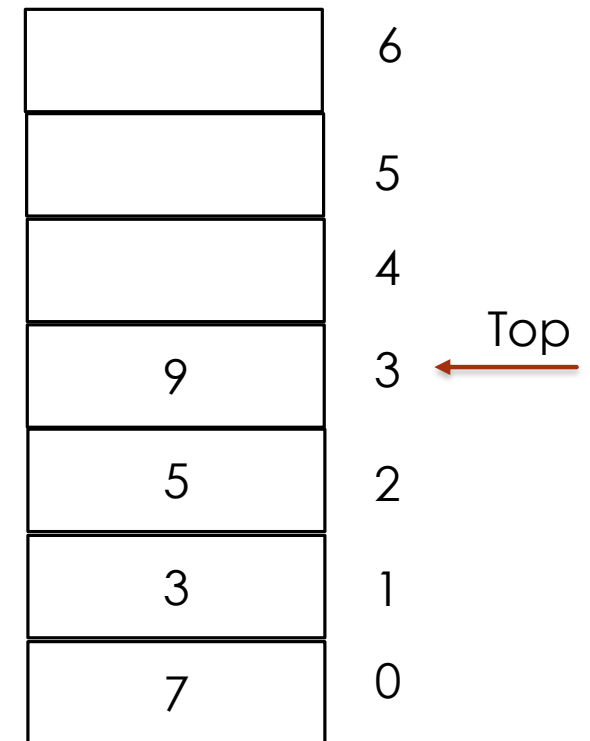


Stack Operations - Example



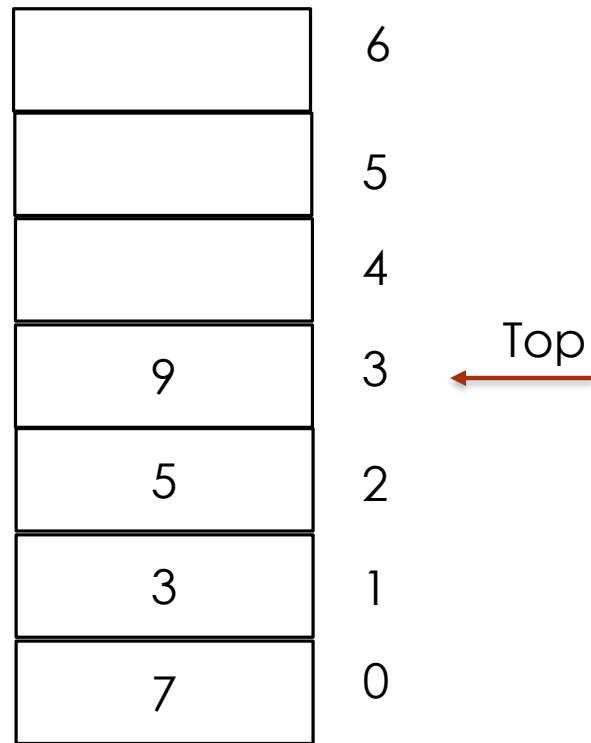
Stack S, N=7

```
Push (9)  
{  
  if top==N-1 Then  
    “Overflow”  
  else  
    Top=Top+1  
    S[top]=9  
}
```



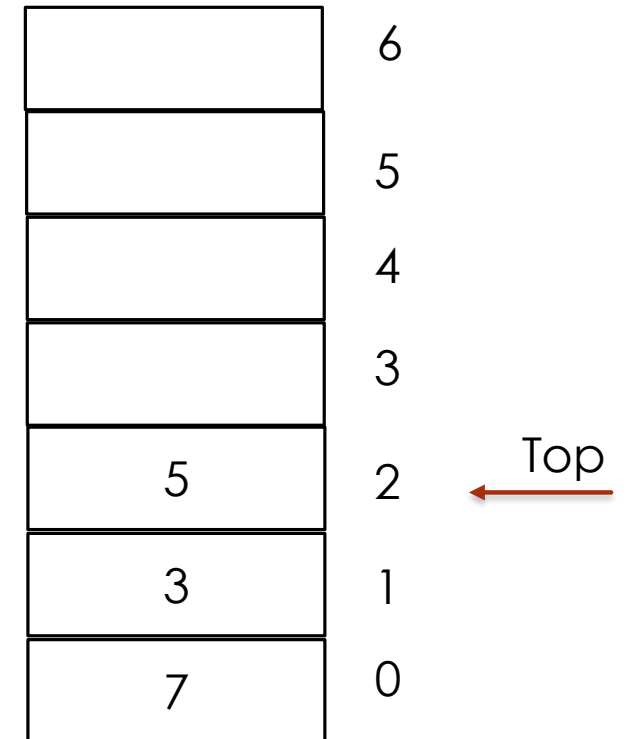
Stack S, N=7

Stack Operations - Example



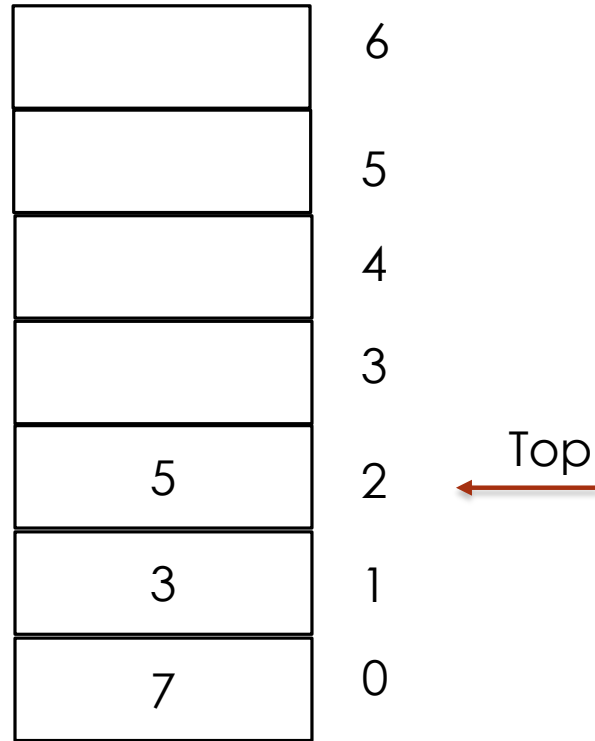
Stack S, N=7

```
Pop ()  
{  
    if isEmpty() then  
        "Underflow"  
    else  
        Top=Top-1  
        return S[top+1]  
}
```



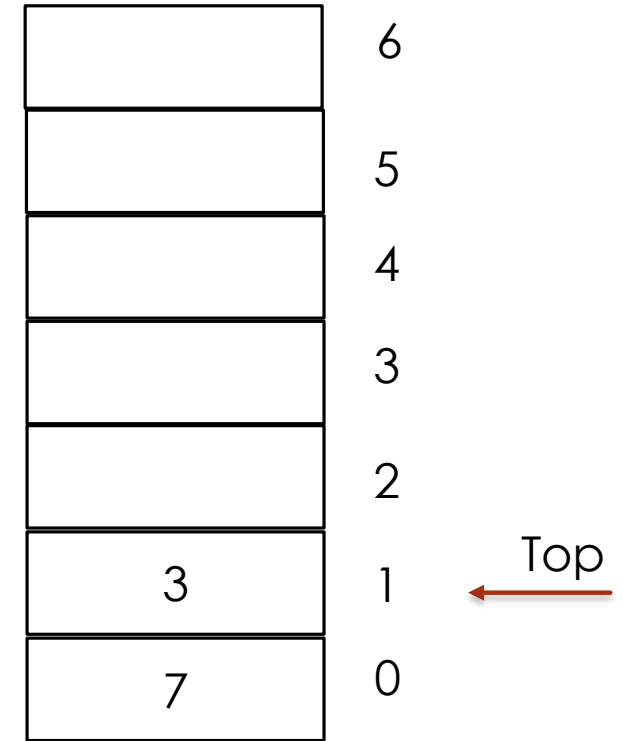
Stack S, N=7

Stack Operations - Example



Stack S, N=6

```
Pop ()  
{  
    if isEmpty() then  
        "Underflow"  
    else  
        Top--  
        return S[top+1]  
}
```



Stack S, N=6



Arithmetic Expression

Arithmetic Expressions

- An **arithmetic expression** is an expression that results in a numeric value.
- It is a correct combination of numbers, operators, parenthesis, and variables.
- Expressions are usually represented in what is known as **Infix notation**, in which each operator is written between two operands
- **Example: $A + B$**
 - **A** and **B** are called **Operands**
 - **+** is called the **operator**

Arithmetic Expressions

➤ Infix form

- **Need precedence rules**
- **May use parentheses.**

➤ Example: $2+4*3$

What is the result?

- Apply precedence rules (***** has higher precedence than **+**)
- We may **use parentheses** rules $(2+4)*3$ or $2+(4*3)$

Rules of Precedence for Arithmetic Operators

Operator	Rule of Precedence
\wedge	Exponentiation (\wedge) is performed first
$*$ /	Multiplication ($*$) and division ($/$) are performed following exponentiation.
$+$ -	Addition ($+$) and subtraction ($-$) are performed last.

- Use parentheses to override precedence rules



There are two more forms for representing an arithmetic expressions in which they do not need **precedence rules or parentheses:**

- **postfix**
- **prefix**

Arithmetic Expressions

- **Postfix form:** Refers to the notation in which the **operator** symbol is placed **after** its two operands
 - Operator appears **after** the operands
 - Infix: $(4+3)*5 \rightarrow$ Postfix: $4\ 3\ +\ 5\ *$
 - Infix: $4+(3*5) \rightarrow$ Postfix: $4\ 3\ 5\ *\ +$
 - **No precedence rules or parentheses!**
- **Prefix Form:** Refers to the notation in which the **operator** symbol is placed **before** its two operands.
 - Operator appears **before** the operands
 - Infix: $(4+3)*5 \rightarrow$ Prefix: $*+4\ 3\ 5$
 - Infix: $4+(3*5) \rightarrow$ Prefix: $+4\ *3\ 5$
 - **No precedence rules or parentheses!**
- **Two Questions:**
 - **How to convert an infix form to postfix and prefix forms.**
 - **How to evaluate an expression given in postfix and prefix forms.**



Stack Applications

Arithmetic Expression

- Conversions
- Evaluations

Example: Infix to Postfix

➤ Example 1 :

$A + B * C + D$

⇒ $A+BC^*+D$

⇒ ABC^*++D

⇒ ABC^*+D+

➤ Example 2:

$A * B + C * D$

⇒ AB^*+C*D

⇒ AB^*+CD^*

⇒ AB^*CD^*+

➤ Example 3:

$A + B * C - D / E * F$

⇒ $A + BC^* - D / E * F$

⇒ $A + BC^* - DE / * F$

⇒ $A + BC^* - DE / F *$

⇒ $A BC^* + - DE / F *$

⇒ $A BC^* + DE / F * -$

➤ Example 4:

$(A+B) * (C+D)$

⇒ $(A B+) * (C+D)$

⇒ $(A B+) * (CD+)$

⇒ $(AB) + (CD+) *$

⇒ $AB + CD+ *$

Example: Infix to Prefix

Example1 :

$A + B * C + D$

$\Rightarrow A+*BC+D$

$\Rightarrow A++*BCD$

$\Rightarrow +A+*BCD$

Example2:

$A * B + C * D$

$\Rightarrow A*B+*CD$

$\Rightarrow *AB+*CD$

$\Rightarrow +*AB*CD$

Example3:

$A +B*C-D/E*F$

$\Rightarrow A +B*C-D/*EF$

$\Rightarrow A +B*C-/D*EF$

$\Rightarrow A +*BC-/D*EF$

$\Rightarrow A +-*BC/D*EF$

$\Rightarrow +A -*BC/D*EF$

Example4:

$(A +B)*(C+D)$

$\Rightarrow (A+ B)*(+CD)$

$\Rightarrow (+A B)*(+CD)$

$\Rightarrow *(+A B)(+CD)$

$\Rightarrow *+A B+CD$

Infix to Postfix Algorithm

While (we have not reached the end of infix expression) // Read from left to right.

If (an **operand** is found) **then**

Add it to **Postfix**

If (a **left parenthesis** '(' is found) **then**

Push it onto the **stack**

If (a **right parenthesis** ')' is found) **then**

While (the stack is not empty AND the top item is not a left parenthesis)

Pop the **stack** and add the popped value to **Postfix**

End-While

Pop the left parenthesis from the **stack** and discard it

If (an **operator** is found) **then**

If (the **stack** is empty or if the top element is a left parenthesis) **then**

Push the operator onto the **stack**

Infix to Postfix Algorithm

Else

While (the **stack** is not empty AND the top of the stack is not a left parenthesis
AND precedence of the **operator** \leq precedence of the **top** of the **stack**)

Pop the **stack** and add the top value to **Postfix**

End-While

Push the latest operator onto the stack

End-While

While (the stack is not empty)

Pop the **stack** and add the popped value to **Postfix**

End-While

While (we have not reached the end of infix expression)

If (an **operand** is found) **then**

Add it to **Postfix**

If (a **left parenthesis** '(' is found) **then**

Push it onto the **stack**

If (a **right parenthesis** ')' is found) **then**

While (the **stack** is not empty AND the top item is not a left parenthesis)

Pop the **stack** and add the popped value to **Postfix**

End-While

Pop the left parenthesis from the **stack** and discard it

If (an **operator** is found) **then**

If (the **stack** is empty or if the top element is a left parenthesis) **then**

Push the operator onto the **stack**

Else

While (the **stack** is not empty AND the top of the stack is not a left parenthesis AND precedence of the **operator** \leq precedence of the **top** of the **stack**)

Pop the **stack** and add the top value to **Postfix**

End-While

Push the latest operator onto the stack

End-While

While (the stack is not empty)

Pop the **stack** and add the popped value to **Postfix**

End-While

Infix to Postfix Algorithm

Infix to Postfix Algorithm Example

➔ Infix Form: $(A+B*C-D)/(E*F)$

Token	Stack	Postfix
((
A	(A
+	(+	
B	(+	AB
*	(+*	
C	(+*	ABC
-	(-	ABC*+
D	(-	ABC*+D
)		ABC*+D-
/	/	
(/(
E	/(ABC*+D-E
*	/(*	
F	/(*	ABC*+D-EF
)	/	ABC*+D-EF*
		ABC*+D-EF*/



- How to convert infix to prefix?
- What is the algorithm of converting an infix to prefix?

Hint: update the little things from the algorithm of converting infix to postfix.

Evaluating a postfix expression Algorithm

While (we have not reached the end of expression) // *Read from left to right.*

If an **operand** is found **then**

push it onto the **stack**

If an **operator** is found **then**

// Pop Twice

A=Pop()

B=Pop()

Evaluate B **operator** A using the operator just found.

Push the resulting value onto the **stack**.

End-While

Pop the stack (this is the final value)

Evaluating a postfix expression - Example

➤ Postfix: $244^*+6-23^*/$

➤ Infix: $(2+4^*4-6)/(2^*3)=2$

Token	Stack
2	2
4	2 4
4	2 4 4
*	2 16
+	18
6	18 6
-	12
2	12 2
3	12 2 3
*	12 6
/	2



How to evaluate a prefix expression?

Hint: update the little things from the algorithm of converting infix to postfix.



Stack Implementation

- ▶ **Array:** We will use this first.
- ▶ **Linked Lists:** Later to be implemented with list.



Lab Assignment

- ▶ Implement the Stack in C++ using OOP.



Exercises

Exercises

- A linear list of elements in which deletion and insertion can be done from one side is known as a?
 - a) Queue.
 - b) Stack.
 - c) Tree.
 - d) Linked list.

- A Stack follows
 - a) FIFO (First In First Out) principle.
 - b) LIFO (Last In First Out) principle.
 - c) Ordered array.
 - d) Linear tree.

Exercises

- Convert the following infix expression to postfix expressions **using Stack** data structure.
 - $(5 * (((9 + 8) * (4 * 6)) + 7))$
 - $6 * (5 + (2 + 3) * 8 + 3)$
- Convert the following infix expression to prefix expressions **using Stack** data structure.
 - $a + b * c + (d * e + f) * g$
- For each of the of the following postfix expressions, find the infix.
 - $6\ 5\ 2\ 3\ +\ 8\ *\ +\ 3\ +\ *$
 - $a\ b\ c\ *\ +\ d\ e\ *\ f\ +\ g\ *\ +$
- Evaluate the following postfix expression
 - $6\ 2\ 5\ 3\ +\ 4\ *\ +\ 3\ +\ *$